Lecture #9: Nondeterministic Time — Speedup, Emulation, and a Nondeterministic Time Hierarchy Theorem Exercises and Review

Additional Exercises

 As noted in the additional exercises for Lecture #8, some authors define nondeterministic time complexity classes using *nondeterministic nondeterministic Turing machines* that, essentially, limit nondeterministic steps to answering Yes-No questions.

Explain why you would have trouble proving a "Nondeterministic Linear Speedup Theorem" (and, indeed, why such a result might not even be correct) when restricted nondeterministic Turing machines are used in this way.

2. Use the Nondeterministic Time Hierarchy Theorem to prove that

 $\mathsf{NTIME}(n) \subsetneq \mathsf{NTIME}(\lceil n\sqrt{n} \rceil) \subsetneq \mathsf{NTIME}(n^2).$

Questions for Review

- 1. State the *Nondeterministic Linear Speedup Theorem*. How does it help to explain why nondeterministic time complexity classes are defined in the way that they are?
- 2. Briefly describe how a nondeterministic *k*-tape Turing machine can be simulated using a 3-tape nondeterministic Turing machine. How are the number of steps used by the *k*-tape machine, and the number of steps used by the 3-tape machine simulating it, related?

- 3. Briefly describe each of the following language and say what its known about its complexity (or, possibly, decidability).
 - (a) *L*_{NTM}
 - (b) $L_{\text{NTM+I}}$
 - (c) $A_{\text{NTM+I}}$
 - (d) $L_{\text{NTM+I+Time}}$
 - (e) $A_{\text{NTM+I+Time}}$
- 4. State the Nondeterministic Time Hierarchy Theorem. Described what it is used for, and summary the steps you must carry out in order to use it.