

- Mandatory: Chapter 3 - Sections 3.3 \& 3.4


At the end of this section, you will be able to:

1. Understand (once more) how graphs are used to represent data
2. Learn the graph coloring algorithm
3.Apply graph coloring to solve for various scheduling problems


Based on Aho, Hopcroft, and Ullman, Data Structures and Algorithms, Addison-Wesley

## Simple Intersection

- Design a traffic light for a complex intersection
- Identify allowed turns
- Going straight is a turn!
- Group turns so that permitted simultaneous turns are in the same group
- Find the smallest number of groups This means a traffic light with the smallest number of phases


Based on Aho, Hopcroft, and Ullman, Data Structures and Algorithms, Addison-Wesley

## Simple Intersection

- From A to B (denoted AB)
- From A to C (AC)
- From A to D (AD)
- BA
- BC
- BD
- DA
- DB
- DC

- AC
- AD
- BA



## Example Simultaneous Turns

- AC conflicts with
- DB
- BD
-BC
- DC

- Build a graph model
- Vertices are turns
- An edge connects conflicting turns

- Color the vertices of the graph
- No two adjacent vertices can have the same color
Adjacent: and edge connects them
- Each color represents a group of turns that can be simultaneous
- Note: more than one solution
- Repeat the following two steps until the graph is colored

1. Select an uncolored vertex and color it with a new color, C
2. For each uncolored vertex, Determine if it has an edge with a vertex that is colored with color C

- If not, color it with color C
- If yes, skip it
- You are to schedule final exams so that a student will not have two final exams scheduled at the same time.
- Although you could schedule each exam in it's own individual time slot (i.e., no two exams run simultaneously) this would be highly inefficient.
- Another constraint is to schedule exams with the minimum number of time slots.
- Examinations for a lecture will be represented with vertices.
- A pair of vertices will be connected if there is at least one student who is registered in both.


## Example 1: Final Exams

- Lecture A:

Student 1, student 2

- Lecture B:

Student 2, student 4, student 6

- Lecture C:

Student 1, student 3, student 5

- Lecture D:
- Student 4, student 5

's Extra, Graph Coloring Example 1:
Final Exams (2)
Peeking into Computer Science
- Lecture A:

Student 1, student 2, student 4

- Lecture B:

Student 2, student 4, student 6

- Lecture C:

Student 1, student 3, student 5

- Lecture D:
- Student 4, student 5



## JT's Extra, Graph Coloring Example 2:

Final Exams
Peeking into Computer Science



## Graph Model of Conflicting Jurns




From Aho, Hopcroft, and Ullman, Data Structures and Algorithms, Addison-Wesley


## Trees

A popular special case of graphs

At the end of this section, the student will be able to:
1.Identify when a graph is a tree
2.Use tree terminology
3.Identify different applications of trees
4.Define binary trees

Objectives

- Trees are a special case of graphs
- A tree is a (directed) graph with the properties:
1.There is a designated vertex, called the root of the tree
2.There is a unique directed path from the root to every other vertex in the tree
- Grow downwards!



## Example Тree



- Root
- Child
- Parent
- Ancestor
- Descendant
- Distance
- Siblings
- Leaf


- A is the parent of B, C, D.
-B, C, D are the children of $A$.
- B,C,D,E are descendants of $A$.
- A is ancestor of B, C, D, E
- All vertices have exactly one parent except for the root (which has none).
- Siblings have the same parent.
- B,C,D are siblings
- E, F are siblings
- A, G have no siblings
- Leaves are vertices with no children
- B, D, E are leaves
- "Bottom level"

- Moving up from a node to the parent is one 'step'.
- Moving up to the parent of the parent is another 'step'.



## Distance






- A binary tree is a tree with the following properties:

1. The edges are labeled with the labels left or right
2. Every vertex has at most two children; if both children exist, then one edge must be labeled with left and the other right.


- Drop the labels, if the tree is drawn with no confusion

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Figure 4.14 Expression tree for $(a+b * c)+((d * e+f) * g)_{1}$


1 From: "Data Structures and Algorithm Analysis in C++", Weiss M.A.

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